

Lecture 6: Message Passing Interface

- Introduction
- The basics of MPI
- Some simple problems
- More advanced functions of MPI
- A few more examples

When is Parallel Implementation Useful

- In general it is useful for Large problems
- Problems suitable for parallelisation, i.e. you know what speed-up to expect
- You need to be able to recognise them
- Three types of problems are suitable:
 - Parallel Problems
 - Regular and Synchronous Problems
 - Irregular and/or Asynchronous Problems

When is Parallel Implementation Useful: Type I

- Parallel problems:
 - The problem can be broken down into subparts
 - Each subpart is independent of the others
 - No communication is required, except to split up the problem and combine the final results
 - Linear speed-up can be expected
- Example of this is: Monte-Carlo simulations

When is Parallel Implementation Useful:

Type II

- Regular and Synchronous Problems:
 - Same instruction set (regular algorithm) applied to all data
 - Synchronous communication (or close to): each processor finishes its task at the same time
 - Local (neighbour to neighbour) and collective (combine final results) communication
- Speed-up based on the computation:communication ratio
- If it is large, expect good speed-up for local communications & ok speed-up for non-local communications
- Ex: Fast Fourier transforms (synchronous), matrix-vector products, sorting (loosely synch.)

When is Parallel Implementation Useful: Type III

- Irregular and/or Asynchronous Problems:
 - Irregular algorithm which cannot be implemented efficiently except with message passing and high communication overhead
 - Communication is usually asynchronous and requires careful coding and load balancing
 - Often dynamic repartitioning of data between processors is required
 - Speed-up is difficult to predict; if the problem can be split up into regular and irregular parts, this makes things easier
- Ex: Melting ice problem (or any moving boundary simulation)

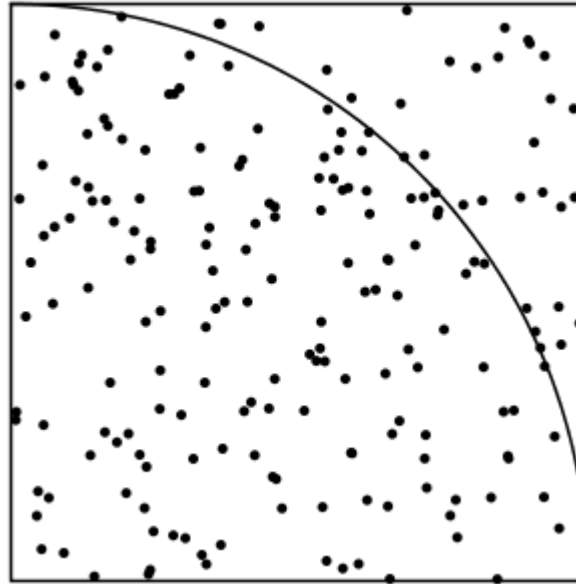
Example 1: matrix-vector product

$$\begin{pmatrix} a_{11} & a_{12} & a_{13} & a_{14} \\ a_{21} & a_{22} & a_{23} & a_{24} \\ a_{31} & a_{32} & a_{33} & a_{34} \\ a_{41} & a_{42} & a_{43} & a_{44} \end{pmatrix} \times \begin{pmatrix} b_1 \\ b_2 \\ b_3 \\ b_4 \end{pmatrix} = \begin{pmatrix} c_1 \\ c_2 \\ c_3 \\ c_4 \end{pmatrix}$$

$$\text{with} \begin{cases} c_1 = a_{11} \times b_1 + a_{12} \times b_2 + a_{13} \times b_3 + a_{14} \times b_4 \\ c_2 = a_{21} \times b_1 + a_{22} \times b_2 + a_{23} \times b_3 + a_{24} \times b_4 \\ c_3 = a_{31} \times b_1 + a_{32} \times b_2 + a_{33} \times b_3 + a_{34} \times b_4 \\ c_4 = a_{41} \times b_1 + a_{42} \times b_2 + a_{43} \times b_3 + a_{44} \times b_4 \end{cases}$$

- A parallel approach:
 - Each element of vector c depends on vector b and only one line of \mathbf{A}
 - Each element of c can be calculated independently from the others
 - Communication only needed to split up the problem and combine the final results
- \Rightarrow a linear speed-up can be expected for large matrices

Example 2: : Monte-Carlo calculation of Pi



Quadrant of a Unit Circle with
Random distribution of points


- $\pi = 3.14159\dots = \text{area of a circle of radius } 1$
- $\pi/4 \approx \text{fraction of the points within the circle quadrant}$
- The more points, the more accurate the value for π is

Example 2: Monte-Carlo calculation of π (Cont'd)

- A parallel approach:
 - Each point is randomly placed within the square
 - The position of each point is independent of the position of the others
 - We can split up the problem by letting each node randomly place a given number of points
 - Communication is only needed to specify the number of points and combine final results
- \Rightarrow a linear speed-up can be expected, allowing for a larger number of points and therefore a greater accuracy in the estimation of π .

Example 3: A More Real Problem

- After each move, the chess software must find the best move within a set
=> This set is large, but finite
- Each move from this set can be evaluated independently & the set can be partitioned
- Communication is only needed to split up the problem and combine the final results
- => A linear speed-up can be expected
- => This means that, in a reasonable time, moves can be studied more thoroughly
- => This depth of evaluation is what makes the software more competitive

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Some background on MPI

- Developed by MPI forum (made up of Industry, Academia and Govt.)
- They established a standardised Message-Passing Interface (MPI-1) in 1994
- It was intended as an interface to both C and FORTRAN.
- C++ bindings were deprecated in MPI-2. Some Java bindings exist but are not standard yet.
- Aim was to provide a specification which can be implemented on any parallel computer or cluster; hence portability of code was a big aim.

Advantages of MPI

- + Portable, hence protection of software investment
- + A standard, agreed by everybody
- + Designed using optimal features of existing message-passing libraries
- + “Kitchen-sink” functionality, very rich environment (129 functions)
- + Implementations for F77, C and C++ are freely downloadable

.....& It's Disadvantages

- “Kitchen-sink” functionality, makes it hard to learn all (unnecessary: a bare dozen are needed in most cases)
- Implementations on shared-memory machines is often quite poor, and does not suit the programming model
- Has rivals in other message-passing libraries (e.g. PVM)

MPI Preliminaries...

- MPI provides support for:
 - Point-to-point & collective (i.e. group) communications
 - Inquiry routines to query the environment (how many nodes are there, which node number am I, etc.)
 - Constants and data-types
- We will start with the basics: initialising MPI, and using point-to-point communication

MPI Preliminaries... (Cont'd)

- Naming convention
 - All MPI identifiers are prefixed by **'MPI_'**.
 - C routines contain lower case (i.e. **'MPI_Init'**),
 - Constants are all in upper case (e.g. **'MPI_FLOAT'** is an MPI C data-type).
 - C routines are actually integer functions which return a status code (you are strongly advised to check these for errors!).
- Running MPI
 - Number of processors used is specified in the command line, when running the MPI loader that loads the MPI program onto the processors, to avoid hard-coding this into the program
 - e.g. **mpirun -np N exec**

MPI Preliminaries... (Cont'd)

- Writing a program using MPI: what is parallel, what is not
 - Only one program is written. By default, every line of the code is executed by each node running the program.
 - E.g. if the code contains `int result=0`, each node will locally create a variable and assign the value.
- When a section of the code needs to be executed by only a subset of nodes, it has to be explicitly specified.
- E.g., providing that we are using 8 nodes, and that **MyID** is a variable storing the rank of the node (from 0 to 7, we will see how to get it later), this section of code assigns to **result** to zero for the first half of them, and 1 for the second.

```
int result;  
    if(MyID < 4) result = 0;  
else result = 1;
```

Common MPI Routines


- MPI has a ‘kitchen sink’ approach of 129 different routines
- Most basic programs can get away with using six.
- As usual use #include “mpi.h” in C.

<code>MPI_Init</code>	Initialise MPI computation
<code>MPI_Finalize</code>	Terminate MPI computation
<code>MPI_Comm_size</code>	Determine number of processes
<code>MPI_Comm_rank</code>	Determine my process number
<code>MPI_Send, MPI_Isend</code>	Blocking, non-blocking send
<code>MPI_Recv, MPI_Irecv</code>	Blocking, non-blocking

Common MPI Routines (Cont'd):

MPI Initialisation, Finalization

- In all MPI-written programs, MPI must be initialised before use, and finalised at the end.
- All MPI-related commands and types must be handled within this section of code:

	<code>MPI_Init</code>	Initialise MPI computation
	<code>⋮</code>	
	<code>MPI_Finalize</code>	Terminate MPI computation

- `MPI_Init` takes two parameters as input (`argc` and `argv`),
 - It is used to start the MPI environment, create the default communicator (more later) and assign a rank to each node.
- `MPI_Finalize` cleans up all MPI state. Once this routine is called, no MPI routine (even `MPI_INIT`) may be called.
- The user must ensure that all pending communications involving a process completes before the process calls `MPI_Finalize`.

Common MPI Routines (Cont'd):

Basic Inquiry Routines

- At various stages in a parallel-implemented function, it may be useful to know how many nodes the program is using, or what the rank of the current node is.
- The `MPI_Comm_size` function returns the number of processes/nodes as an integer, taking only one parameter, a communicator.
- In most cases you will only use the default Communicator: `MPI_COMM_WORLD`.
- The `MPI_Comm_rank` function is used to determine what the rank of the current process/node on a particular communicator.
- E.g. if there are two communicators, it is possible, and quite usual, that the ranks of the same node would differ.
- Again, in most cases, this function will only be used with the default communicator as an input (`MPI_COMM_WORLD`), and it will return (as an integer) the rank of the node on that communicator.

Common MPI Routines (Cont'd): Point-to-Point communications in MPI

- This involves communication between two processors, one sending, and the other receiving.
- Certain information is required to specify the message:
 - Identification of sender processor
 - Identification of destination/receiving processor
 - Type of data (`MPI_INT`, `MPI_FLOAT` etc)
 - Number of data elements to send (i.e. array/vector info)
 - Where the data to be sent is in memory (pointer)
 - Where the received data should be stored in (pointer)

Common MPI Routines (Cont'd):

Sending data `MPI_Send`, `MPI_Isend`

- `MPI_Send` is used to perform a blocking send, (i.e. process waits for the communication to finish before going to the next command).

```
int MPI_Send(void *buf, int count, MPI_Datatype datatype, int dest, int tag, MPI_Comm comm)
```

- This function takes six parameters:
 - the location of the data to be sent i.e. a pointer (*input parameter*)
 - the number of data elements to be sent (*input parameter*)
 - the type of data e.g. `MPI_INT`, `MPI_FLOAT`, etc. (*input parameter*)
 - the rank of the receiving/destination node (*input parameter*)
 - a tag for identification of the communication (*input parameter*)
 - the communicator to be used for transmission (*input parameter*)
- `MPI_Isend` is non-blocking, so an additional parameter, to allow for verification of communication success is needed.
- It is a pointer to an element of type `MPI_Request`.

Common MPI Routines (Cont'd):

Receiving data `MPI_Recv`, `MPI_Irecv`

- `MPI_Recv` is used to perform a blocking receive, (i.e. process waits for the communication to finish before going to the next command).

```
int MPI_Recv(void *buf, int count, MPI_Datatype datatype, int source, int tag, MPI_Comm comm, MPI_Status *status);
```

- This function takes seven parameters:
 - the location of the receive buffer i.e. a pointer (*output parameter*)
 - the max number of data elements to be received (*input parameter*)
 - the type of data e.g. `MPI_INT`, `MPI_FLOAT`, etc. (*input parameter*)
 - the rank of the source/sending node (*input parameter*)
 - a tag for identification of the communication (*input parameter*)
 - the communicator to be used for transmission (*input parameter*)
 - a pointer to a structure of type `MPI_Status`, contains source processor's rank, communication tag, and error status (*output parameter*)
- For the non-blocking `MPI_Irecv`, `MPI_Request` replaces `MPI_Status`.

A first MPI example: Hello World.

```
#include <mpi.h>
int main(int argc, char *argv[]) {
    int myid, size;
    MPI_Init(&argc, &argv);
    MPI_Comm_rank(MPI_COMM_WORLD, &myid);
    MPI_Comm_size(MPI_COMM_WORLD, &size);
    printf("process %d out of %d says Hello\n", myid, size);
    MPI_Finalize();
    return 0;
}
```

First “real” MPI program: Exchanging 2 Values

```
#include <mpi.h>
int main(int argc, char *argv[]) {
    int rank, value, size, length = 1, tag = 1;
    MPI_Status status;
    /* initialize MPI and get own id (rank) */
    MPI_Init(&argc, &argv);
    MPI_Comm_rank(MPI_COMM_WORLD, &myid);
    MPI_Comm_size(MPI_COMM_WORLD, &size);
    if (size!=2) {
        printf("use exactly two processes\n");
        exit(1);
    }
    if (myid == 0) {
        otherid = 1; myvalue = 14;
    }
    else {
        otherid = 0; myvalue = 25;
    }

    printf("process %d sending %d to process %d\n", myid, myvalue, otherid);
    /* Send one integer to the other node (i.e. "otherid") */
    MPI_Send(&myvalue,1,MPI_INT,otherid,tag,MPI_COMM_WORLD);
    /* Receive one integer from any other node */
    MPI_Recv(&othervalue,1,MPI_INT,MPI_ANY_SOURCE,
            MPI_ANY_TAG,MPI_COMM_WORLD, &status);
    printf("process %d received a %d\n", myid, othervalue);
    MPI_Finalize();          /* Terminate MPI */
    return 0;
}
```

Compiling and Running MPI Programmes

- To compile programs using MPI, you need an “MPI-enabled” compiler.
- On our cluster, we use `mpicc` to compile C programs containing MPI commands or `mpicxx` for C++.
- Before running an executable using MPI, you need to make sure the "multiprocessing daemon" (MPD) is running.
- It makes the workstations into "virtual machines" to run MPI programs.
- When you run an MPI program, requests are sent to MPD daemons to start up copies of the program.
- Each copy can then use MPI to communicate with other copies of the same program running in the virtual machine. Just type “mpd &” in the terminal.
- To run the executable, type “`mpirun -np N./executable_file`”, where N is the number to be used to run the program.
- This value is then used in your program by `MPI_Init` to allocate the nodes and create the default communicator.

Example 3: “Ring” Communication

```
#include <mpi.h>
int main(int argc, char *argv[]) {

    int rank, value, size;
    MPI_Status status;
    /* initialize MPI and get own id (rank) */
    MPI_Init(&argc, &argv);
    MPI_Comm_rank(MPI_COMM_WORLD, &rank);
    MPI_Comm_size(MPI_COMM_WORLD, &size);
    do {
        if (rank == 0) {
            scanf("%d", &value);
            /* Master Node sends out the value */
            MPI_Send(&value, 1, MPI_INT, rank + 1, 0, MPI_COMM_WORLD);
        }
        else {
            /* Slave Nodes block on receive the send on the value */
            MPI_Recv(&value, 1, MPI_INT, rank - 1, 0, MPI_COMM_WORLD, &status);
            if (rank < size - 1) {
                MPI_Send(&value, 1, MPI_INT, rank + 1, 0, MPI_COMM_WORLD);
            }
            printf("process %d got %d\n", rank, value);
        } while (value >= 0);
        /* Terminate MPI */
    }
    MPI_Finalize();
    return 0;
}
```



```

#include <mpi.h>
int main(int argc, char *argv[]) {
    int A[4][4], b[4], c[4], line[4], temp[4], local_value, myid;
    MPI_Init(&argc, &argv); MPI_Comm_rank(MPI_COMM_WORLD, &myid);
    if (myid == 0) {
        for (int i=0; i<4; i++) {
            b[i] = 4 - i;
            for (int j=0; j<4; j++)
                A[i][j] = i + j; /* set some notional values for A, b */
        }
        line[0]=A[0][0]; line[1]=A[0][1];
        line[2]=A[0][2]; line[3]=A[0][3];
    }
    if (myid == 0) {
        for (int i=1; i<4; i++) { /* slaves do most of the multiplication */
            temp[0]=A[i][0]; temp[1] = A[i][1]; temp[2] = A[i][2]; temp[3] = A[i][3];
            MPI_Send( temp, 4, MPI_INT, i, i, MPI_COMM_WORLD);
            MPI_Send( b, 4, MPI_INT, i, i, MPI_COMM_WORLD);
        }
    }
    else {
        MPI_Recv( line, 4, MPI_INT, 0, myid, MPI_COMM_WORLD, MPI_STATUS_IGNORE);
        MPI_Recv( b, 4, MPI_INT, 0, myid, MPI_COMM_WORLD, MPI_STATUS_IGNORE);
        /* master node does its share of multiplication too */
        c[myid] = line[0] * b[0] + line[1] * b[1] + line[2] * b[2] + line[3] * b[3];
        if (myid != 0) {
            MPI_Send(&c[myid], 1, MPI_INT, 0, myid, MPI_COMM_WORLD);
        }
        else {
            for (int i=1; i<4; i++) {
                MPI_Recv( &c[i], 1, MPI_INT, i, i, MPI_COMM_WORLD, MPI_STATUS_IGNORE);
            }
        }
    }
    MPI_Finalize();
    return 0;
}

```

Example 4: Matrix-Vector Product Implementation

Example 5: Pi Calculation Implementation

```
int main(int argc, char *argv[]) {
    MPI_Init(&argc, &argv);
    #define INT_MAX_ 1000000000
    int myid, size, inside=0, outside=0, points=10000;
    double x,y, Pi_comp, Pi_real=3.141592653589793238462643;
    MPI_Comm_rank(MPI_COMM_WORLD, &myid);
    MPI_Comm_size(MPI_COMM_WORLD, &size);
    if (myid == 0) {
        for (int i=1; i<size; i++) /* send out the value of points to all slaves */
            MPI_Send(&points, 1, MPI_INT, i, i, MPI_COMM_WORLD);
    }
    else
        MPI_Recv(&points, 1, MPI_INT, 0, i, MPI_COMM_WORLD, MPI_STATUS_IGNORE);
    rands=new double[2*points];
    for (int i=0; i<2*points; i++){
        rands[i]=random();
        if (rands[i]<=INT_MAX_) i++ /* this random is within range */
    }
    for (int i=0; i<points;i++){
        x=rands[2*i]/INT_MAX_;
        y=rands[2*i+1]/INT_MAX_;
        if((x*x+y*y)<1) inside++ /* point is inside unit circle so incr var inside */
    }
    delete[] rands;
    if (myid == 0) {
        for (int i=1; i<size; i++) {
            int temp; /* master receives all inside values from slaves */
            MPI_Recv(&temp, 1, MPI_INT, i, i, MPI_COMM_WORLD, MPI_STATUS_IGNORE);
            inside+=temp; /* master sums all insides sent to it by slaves */
        }
    }
    else
        MPI_Send(&inside, 1, MPI_INT, 0, i, MPI_COMM_WORLD); /* send inside to master */
    if (myid == 0) {
        Pi_comp = 4 * (double) inside / (double) (size*points);
        cout << "Value obtained: " << Pi_comp << endl << "Pi:" << Pi_real << endl;
    }
    MPI_Finalize(); return 0;
}
```

All nodes do this part

Collective communications in MPI

- Groups are sets of processors that communicate with each other in a certain way.
- Such communications permit a more flexible mapping of the language to the problem (allocation of nodes to subparts of the problem etc).
- MPI implements Groups using data objects called Communicators.
- A special Communicator is defined (called '**MPI_COMM_WORLD**') for the group of all processes.
- Each Group member is identified by a number (its Rank 0..n-1).
- There are three steps to create new communication structures:
 - accessing the group corresponding to MPI_COMM_WORLD,
 - using this group to create sub-groups,
 - allocating new communicators for this group.
- We will see this in more detail in the last examples.

Some Sophisticated MPI Routines

- The advantage of the global communication routines below is that the MPI system can implement them more efficiently than the programmer, involving far less function calls.
- Also the system will have more opportunity to overlap message transfers with internal processing and to exploit parallelism that might be available in the communications network.

MPI_Barrier	Synchronise
MPI_Bcast	Broadcast same data to all procs
MPI_Gather	Get data from all procs
MPI_Scatter	Send different data to all procs
MPI_Reduce	Combine data from all onto one proc
MPI_Allreduce	Combine data from all procs onto all procs

Sophisticated MPI Routines:

`MPI_Barrier`

- **`MPI_Barrier`** is used to synchronise a set of nodes.

`int MPI_Barrier(MPI_Comm comm)`

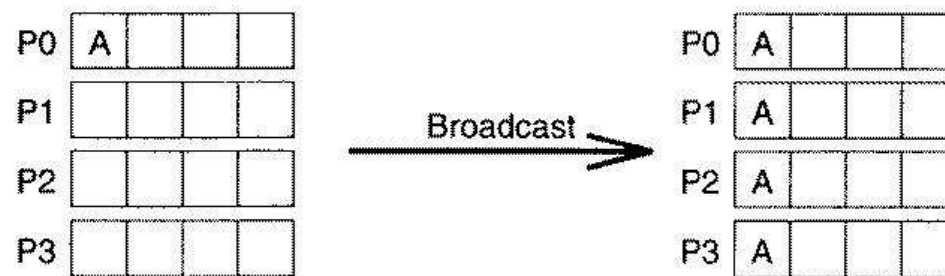
- It blocks the caller until all group members have called it.
- ie call returns at any process only after all group members have entered the call.
- This functions takes only parameter, the communicator (i.e. group of nodes) to be synchronised.
- As we previously saw with other functions, it will most of the times be used with the default communicator, **`MPI_COMM_WORLD`**.

Sophisticated MPI Routines: `MPI_Bcast`

- `MPI_Bcast` used to send data from one node to all the others in one single command.

```
int MPI_Bcast( void *buffer, int count, MPI_Datatype  
datatype, int root, MPI_Comm comm )
```

- This function takes five parameters: -
 - location of data to be sent i.e. a pointer (*input/output* parameter)
 - number of data elements to be sent (*input* parameter)
 - type of data (*input* parameter)
 - rank of the broadcast node (*input* parameter)
 - communicator to be used (*input* parameter)

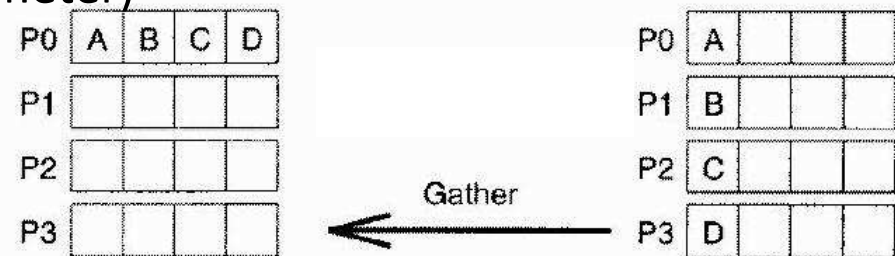


Sophisticated MPI Routines: `MPI_Gather`

- `MPI_Gather` is used to gather on a single node data scattered over a group of nodes.

```
int MPI_Gather(void *sendbuf, int sendcnt, MPI_Datatype sendtype,  
void *recvbuf, int recvcnt, MPI_Datatype recvtype, int root,  
MPI_Comm comm)
```

- This function takes eight parameters: -
 - location of data to be sent i.e. a pointer (*input parameter*)
 - number of data elements to be sent (*input parameter*)
 - type of data to be sent (*input parameter*)
 - location of the receive buffer i.e. a pointer (*output parameter*)
 - number of elements to be received (*input parameter*)
 - type of data to be received (*input parameter*)
 - rank of the sending node (*input parameter*)
 - communicator to be used for transmission. (*input parameter*)

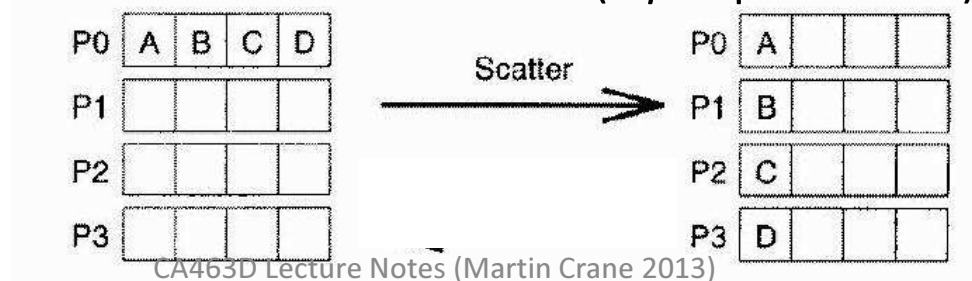


Sophisticated MPI Routines: **MPI_Scatter**

- **MPI_Scatter** used to scatter data from single node to a group

```
int MPI_Scatter(void *sendbuf, int sendcnt, MPI_Datatype  
sendtype, void *recvbuf, int recvcnt, MPI_Datatype recvtype, int  
root, MPI_Comm comm)
```

- This function takes eight parameters: -
 - location of data to be sent i.e. a pointer (*input parameter*)
 - number of data elements to be sent (*input parameter*)
 - type of data to be sent (*input parameter*)
 - location of the receive buffer i.e. a pointer (*output parameter*)
 - number of elements to be received (*input parameter*)
 - type of data to be received (*input parameter*)
 - rank of the sending node (*input parameter*)
 - communicator to be used for transmission. (*input parameter*)



Sophisticated MPI Routines: **`MPI_Reduce`**

- **`MPI_Reduce`** used to reduce values on all nodes of a group to a single value on one node using some reduction operation (sum etc).

```
int MPI_Reduce ( void *sendbuf, void *recvbuf, int count,  
MPI_Datatype datatype, MPI_Op op, MPI_Comm comm )
```

- This function takes six parameters: -
 - location of the data to be sent i.e. a pointer (*input parameter*)
 - location of the receive buffer i.e. a pointer (*output parameter*)
 - number of elements to be sent (*input parameter*)
 - type of data e.g. **`MPI_INT`**, **`MPI_FLOAT`**, etc. (*input parameter*)
 - operation to combine the results e.g. **`MPI_SUM`** (*input parameter*)
 - communicator used for transmission (*input parameter*)



Sophisticated MPI Routines: `MPI_Allreduce`

- `MPI_Allreduce` is used to reduce values on all group nodes to a one value, and send it back to all (i.e. equals `MPI_Reduce`+`MPI_Bcast`)

```
int MPI_Allreduce ( void *sendbuf, void *recvbuf, int count,  
MPI_Datatype datatype, MPI_Op op, MPI_Comm comm )
```

- This function takes six parameters: -
 - location of the data to be sent i.e. a pointer (*input parameter*)
 - location of the receive buffer i.e. a pointer (*output parameter*)
 - number of elements to be sent (*input parameter*)
 - type of data e.g. `MPI_INT`, `MPI_FLOAT`, etc. (*input parameter*)
 - operation to combine the results e.g. `MPI_SUM` (*input parameter*)
 - communicator used for transmission (*input parameter*)



```

#include <mpi.h>
int main(int argc, char *argv[]) {
    int A[4][4], b[4], c[4], line[4], temp[4], local_value, myid;
    MPI_Init(&argc, &argv); MPI_Comm_rank(MPI_COMM_WORLD, &myid);
    if (myid == 0) {
        for (int i=0; i<4; i++) {
            b[i] = 4 - i;
            for (int j=0; j<4; j++)
                A[i][j] = i + j; /* set some notional values for A, b */
        }
        line[0]=A[0][0]; line[1]=A[0][1];
        line[2]=A[0][2]; line[3]=A[0][3];
    }
    MPI_Bcast(b, 4, MPI_INT, 0, MPI_COMM_WORLD);
    if (myid == 0) {
        for (int i=0; i<4; i++) { /* slaves do most of the multiplication */
            temp[0]=A[i][0]; temp[1] = A[i][1]; temp[2] = A[i][2]; temp[3] = A[i][3];
            MPI_Send( temp, 4, MPI_INT, i, i, MPI_COMM_WORLD);
            /* No need to send vector b here */
        }
    }
    else {
        MPI_Recv( line, 4, MPI_INT, 0, myid, MPI_COMM_WORLD, MPI_STATUS_IGNORE);
        /* No need to receive vector b here */
    }
    /* master node does its share of multiplication too */
    c[myid] = line[0] * b[0] + line[1] * b[1] + line[2] * b[2] + line[3] * b[3];
    if (myid != 0) {
        MPI_Send(&c[myid], 1, MPI_INT, 0, myid, MPI_COMM_WORLD);
    }
    else {
        {
            for (int i=1; i<4; i++) {
                MPI_Recv( &c[i], 1, MPI_INT, i, i, MPI_COMM_WORLD, MPI_STATUS_IGNORE);
            }
        }
    }
    MPI_Finalize(); return 0;
}

```

Example 6: A New Matrix-Vector Product

Example 5: A New Pi Calculation Implementation

```
int main(int argc, char *argv[]) {
    MPI_Init(&argc, &argv);
    #define INT_MAX_ 1000000000
    int myid, size, inside=0, outside=0, points=10000;
    double x,y, Pi_comp, Pi_real=3.141592653589793238462643;
    MPI_Comm_rank(MPI_COMM_WORLD, &myid);
    MPI_Comm_size(MPI_COMM_WORLD, &size);
    /* Again send/receive replaced by MPI_Bcast */
    MPI_Bcast(&points,1,MPI_INT, 0, MPI_COMM_WORLD);
    rands=new double[2*points];
    for (int i=0; i<2*points; i++){
        rands[i]=random();
        if (rands[i]<=INT_MAX_) i++ /* this random is within range */
    }
    for (int i=0; i<points;i++){
        x=rands[2*i]/INT_MAX_;
        y=rands[2*i+1]/INT_MAX_;
        if((x*x+y*y)<1) inside++ /* point is inside unit circle so incr var inside */
    }
    delete[] rands;
    if (myid == 0) {
        for (int i=1; i<size; i++) {
            int temp; /* master gets all inside values from slaves */
            MPI_Recv(&temp, 1, MPI_INT, i, i, MPI_COMM_WORLD, MPI_STATUS_IGNORE);
            inside+=temp; /* master sums all insides sent to it by slaves */
        }
    }
    else
        MPI_Send(&inside, 1, MPI_INT, 0, i, MPI_COMM_WORLD); /* send inside to master */
    MPI_Reduce(&inside,&total,1,MPI_INT,MPI_SUM,0, MPI_COMM_WORLD);
    if (myid == 0) {
        Pi_comp = 4 * (double) inside / (double) (size*points);
        cout << "Value obtained: " << Pi_comp << endl << "Pi:" << Pi_real << endl;
    }
    MPI_Finalize(); return 0;
}
```

Using Communicators

- Creating a new group (and communicator) by excluding the first node:

```
#include <mpi.h>
int main(int argc, char *argv[]) {
    :
    :
    MPI_Comm comm_world, comm_worker;
    MPI_Group group_world, group_worker;
    comm_world = MPI_COMM_WORLD;
    MPI_Comm_group(comm_world, &group_world);
    MPI_Group_excl(group_world, 1, 0, &group_worker);
    /* process 0 not member */
    MPI_Comm_create(comm_world, group_worker, &comm_worker);
    :
    :
}
```

- **Warning:**

`MPI_Comm_create()` is a collective operation, so all processes in the old communicator must call it - even those not going in the new communicator.

Example 8: Using Communicators

```
#include <mpi.h>
#include <stdio.h>
#define NPROCS 8
int main(int argc, char *argv[]) {

    int rank, newrank, sendbuf, recvbuf;
    ranks1[4]={0,1,2,3}, ranks2[4]={4,5,6,7};
    MPI_Group orig_group, new_group;
    MPI_Comm new_comm;

    MPI_Init(&argc, &argv);
    MPI_Comm_rank(MPI_COMM_WORLD, &rank);
    sendbuf = rank;

    /* Extract the original group handle */
    MPI_Comm_group(MPI_COMM_WORLD, &orig_group);
    if (rank < NPROCS/2) { /* Split tasks into 2 distinct groups based on rank */
        MPI_Group_incl(orig_group, NPROCS/2, ranks1, &new_group);
    } else {
        MPI_Group_incl(orig_group, NPROCS/2, ranks2, &new_group);
    }
    /* Create new communicator and then perform collective communications */
    MPI_Comm_create(MPI_COMM_WORLD, new_group, &new_comm);
    MPI_Allreduce(&sendbuf, &recvbuf, 1, MPI_INT, MPI_SUM, new_comm);
    MPI_Group_rank (new_group, &new_rank);

    printf("rank= %d newrank= %d recvbuf= %d\n", rank, newrank, recvbuf);

    MPI_Finalize();
}
```

Final Reminder

- MPI programs need specific compilers (e.g. `mpicc`), `MPD` and `mpirun`.
- MPI programs start with `MPI_Init` and finish with `MPI_Finalize`,
- Four functions for point-to-point communication,
- Six more advanced functions, for synchronise, and perform collective communication,
- Nine functions (at least three!) to create new groups and communicators,
- Too many examples to remember everything.